

Algorithms, Probability, and Computing **Fall 09**
Exercise Set 8**Exercise 1**

We consider the following algorithm BETTERRMINCUT for minimum cut, with an input graph G on n vertices and with additional parameters t and T :

BETTERRMINCUT(G, t, T):
 $H \leftarrow \text{RANDOMCONTRACT}(G, t)$
return minimum of
 T runs of BASICMINCUT(H)

RANDOMCONTRACT(G, t):
while $|V(G)| > t$ do
 for random $e \in E(G)$
 $G \leftarrow G/e$
end while
return G

BASICMINCUT(G):
while G has more than 2 vertices do
 pick a random edge e in G
 $G \leftarrow G/e$
end while
return the size of the only cut in G

In the following, you are asked to analyze the behavior of this algorithm for various settings of t and T . The goal is to make really (!) rough calculations to realize what is going on. You need not worry about integer parts, say. For the probability of success of $\text{RANDOMCONTRACT}(G, t)$ use the bound derived in class, which is roughly t^2/n^2 , $t \geq 2$.

- (a) Consider BETTERRMINCUT called with $t = n/2$ and $T = 2$. How do the probability of success and the running time change compared to the basic guessing algorithm BASICMINCUT(G)?
- (b) Consider BETTERRMINCUT called with $t = \sqrt{n}$ and $T = n$. Estimate the probability of success. How many times do we need to repeat this algorithm, in order to make the probability of success at least $\frac{1}{2}$? What is the total running time of these repetitions?
- (c) Consider now $t = n^\alpha$, $T = n^\beta$, where $\alpha \in (0, 1)$ and $\beta > 0$ are constants. Again estimate the success probability and number N of repetitions needed to make the success probability at least $\frac{1}{2}$. What choice of α and β give the best total running time? (If you cannot determine the very best ones, at least give the best ones you can find.)

Exercise 2

Prove that no graph on n vertices has more than $\binom{n}{2}$ minimum cuts.

Hint: Compare $\binom{n}{2}$ with $P_0(n)$, the success probability of BASICMINCUT.

Exercise 3

Suppose we are running the checking algorithm for matrices over $\text{GF}(2)$, i.e. numbers are $\{0, 1\}$ with addition and multiplication mod 2. Show that in one iteration the success probability of detecting an error in the supposed product matrix C is exactly $\frac{1}{2}$, in case matrix C is wrong in exactly one row.