

When is a graph planar?\_\_\_\_\_

**Theorem**(Euler, 1758) If a plane multigraph  $G$  with  $k$  components has  $n$  vertices,  $e$  edges, and  $f$  faces, then

$$n - e + f = 1 + k.$$

**Corollary** If  $G$  is a simple, planar graph with  $n(G) \geq 3$ , then  $e(G) \leq 3n(G) - 6$ .

If also  $G$  is triangle-free, then  $e(G) \leq 2n(G) - 4$ .

**Corollary**  $K_5$  and  $K_{3,3}$  are non-planar.

The **subdivision of edge**  $e = xy$  is the replacement of  $e$  with a new vertex  $z$  and two new edges  $xz$  and  $zy$ . The graph  $H'$  is a **subdivision of  $H$** , if one can obtain  $H'$  from  $H$  by a series of edge subdivisions. Vertices of  $H'$  with degree at least three are called **branch vertices**.

**Theorem**(Kuratowski, 1930) A graph  $G$  is planar **iff**  $G$  does not contain a subdivision of  $K_5$  or  $K_{3,3}$ .

## Kuratowski's Theorem\_\_\_\_\_

**Theorem**(Kuratowski, 1930) A graph  $G$  is planar **iff**  $G$  does not contain a subdivision of  $K_5$  or  $K_{3,3}$ .

*Proof.*

A **Kuratowski subgraph** of  $G$  is a subgraph of  $G$  that is a subdivision of  $K_5$  or  $K_{3,3}$ . A **minimal nonplanar graph** is a nonplanar graph such that every proper subgraph is planar.

A **counterexample** to Kuratowski's Theorem constitutes a **nonplanar** graph that does **not** contain any Kuratowski subgraph.

Kuratowski's Theorem follows from the following Main Lemma and Theorem.

## The spine of the proof\_\_\_\_\_

**Main Lemma.** If  $G$  is a graph with fewest edges among counterexamples, then  $G$  is 3-connected.

**Lemma 1.** Every minimal nonplanar graph is 2-connected.

**Lemma 2.** Let  $S = \{x, y\}$  be a separating set of  $G$ . If  $G$  is a nonplanar graph, then adding the edge  $xy$  to some  $S$ -lobe of  $G$  yields a nonplanar graph.

**Main Theorem.**(Tutte, 1960) If  $G$  is a 3-connected graph with no Kuratowski subgraph, then  $G$  has a convex embedding in the plane with no three vertices on a line.

A **convex embedding** of a graph is a planar embedding in which each face boundary is a convex polygon.

**Lemma 3.** If  $G$  is a 3-connected graph with  $n(G) \geq 5$ , then there is an edge  $e \in E(G)$  such that  $G \cdot e$  is 3-connected.

*Notation:*  $G \cdot e$  denotes the graph obtained from  $G$  after the **contraction** of edge  $e$ .

**Lemma 4.**  $G$  has no Kuratowski subgraph  $\Rightarrow G \cdot e$  has no Kuratowski subgraph.

## Proof of Tutte's Theorem

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**Main Theorem.** (Tutte, 1960) If  $G$  is a 3-connected graph with no Kuratowski subgraph, then  $G$  has a convex embedding in the plane with no three vertices on a line.

*Proof.* Induction on  $n(G)$ .

Base case:  $G$  is 3-connected,  $n(G) = 4 \Rightarrow K_4$ .

Let  $e \in E$  s.t.  $H = G \cdot e$  is 3-connected. (Lemma 3)

Then  $H$  has no Kuratowski subgraph. (Lemma 4)

Induction  $\Rightarrow H$  has a convex embedding in the plane with no three vertices on a line.

Let  $z \in V(H)$  be the contracted  $e$ .

$H - z$  is 2-connected  $\Rightarrow$  boundary of the face containing  $z$  after the deletion of the edges incident to  $z$  is a cycle  $C$ .

Let  $x_0, \dots, x_{k-1}$  be the neighbors of  $x$  on  $C$  in cyclic order. Note that  $|N(x)| \geq 3$  and hence  $k \geq 2$ .

Denote by  $\langle x_i, x_{i+1} \rangle$  the portion of  $C$  from  $x_i$  to  $x_{i+1}$  (including endpoints; indices taken modulo  $k$ ).

Let  $N_x = N(x) \setminus \{y\}$  and  $N_y = N(y) \setminus \{x\}$ .

**Case 1.**  $|N_x \cap N_y| \geq 3$ .

Let  $u, v, w \in N_x \cap N_y$ . Then  $x, y, u, v, w$  are the branch vertices of a  $K_5$ -subdivision in  $G$ .

**Case 2.**  $|N_x \cap N_y| \leq 2$ .

Since  $|N_x \cup N_y| \geq 3$ , there is w.l.o.g. a vertex  $u \in N_y \setminus N_x$ . Let  $i$  be such that  $u$  is on  $\langle x_i, x_{i+1} \rangle$ .

**Case 2a.**  $N_y$  is contained in  $\langle x_i, x_{i+1} \rangle$ .

Then there is an appropriate embedding of  $G$ : Placing  $x$  in place of  $z$  and  $y$  sufficiently close to  $x$  maintains convexity. (No three vertices are collinear;  $|N(x)|, |N(y)| \geq 3$ .)

**Case 2b.** For every  $i$  there is a vertex in  $N_y$  that is not in  $\langle x_i, x_{i+1} \rangle$ .

Then there must be a  $v \in N_y$  that is not on  $\langle x_i, x_{i+1} \rangle$  and  $x, y, x_i, x_{i+1}, u, v$  are the branch vertices of a  $K_{3,3}$ -subdivision in  $G$ .

## Proof of the Lemmas

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**Lemma 3.**  $G$  is 3-connected,  $n(G) \geq 5 \Rightarrow$  there is an edge  $e \in E(G)$  such that  $G \cdot e$  is 3-connected.

*Proof.* Suppose  $G$  is 3-connected and for every  $e \in E$ ,  $G \cdot e$  is NOT 3-connected.

For edge  $e = xy$ , the vertex  $z$  is the mate of  $xy$  if  $\{x, y, z\}$  is a cut in  $G$ .

Choose  $e = xy$  and their mate  $z$  such that  $G - \{x, y, z\}$  has a component  $H$ , whose order is as large as possible.

Let  $H'$  be another component of  $G - \{x, y, z\}$  and let  $u \in V(H')$  be the neighbor of  $z$  (There IS one!). Let  $v$  be the mate of  $uz$ .

$V(H) \cup \{x, y\} \setminus \{v\}$  is connected in  $G - \{z, u, v\}$  contradicting the maximality of  $H$ .  $\square$

**Lemma 4.**  $G$  has **no Kuratowski** subgraph  $\Rightarrow G \cdot e$  has **no Kuratowski** subgraph.

*Proof.* Suppose  $G \cdot e$  contains a Kuratowski subgraph  $H$ . Then

- $z \in V(H)$
- $z$  is a branchvertex of  $H$
- $|N_H(z)| = 4$  and  $|N_H(z) \cap N_G(x)|, |N_H(z) \cap N_G(y)| \geq 2$

Then  $H$  is the subdivision of  $K_5 \Rightarrow G$  contains a subdivision of  $K_{3,3}$ , a **contradiction**. □

## Minors

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$K_7$  is a toroidal graph (it is embeddable on the torus),  $K_8$  is not. What else is not? For the torus there is NO equivalent version of Kuratowski's characterization with a finite number of forbidden subdivisions. Any such characterization must lead to an infinite list.

A weaker concept: Minors.

Graph  $G$  is called a **minor of graph  $H$**  if  $G$  can be obtained from  $H$  by a series of edge deletions and edge contractions. Graph  $H$  is also called a  **$G$ -minor**

*Example:*  $K_5$  is a minor of the Petersen graph  $P$ , but  $P$  does not contain a  $K_5$ -subdivision.

## The Graph Minor Theorem\_\_\_\_\_

**Theorem.** (Robertson and Seymour, 1985-2005) In any infinite list of graphs, some graph is a minor of another.

*Proof:* more than 500 pages in 20 papers.

**Corollary** For any graph property that is closed under taking minors, there exists **finitely many** minimal **forbidden** minors.

*Homework.* Wagner's Theorem. Every nonplanar graph contains either a  $K_5$  or  $K_{3,3}$ -minor.

For embeddability on the **projective plane**, it is known that there are **35** minimal forbidden minors. For embeddability on the **torus**, we don't know the exact number of minimal forbidden minors; there are **more than 800 known**.