

Solution to Exercise 8. Types of intermediate results are either `int` or `unsigned int` (indicated with `u`).

- a) $-2-4*3$ is parenthesized as $(-2)-(4*3)$ since unary minus binds more strongly than binary minus, and since multiplicative operators bind more strongly than additive ones. Evaluation: $(-2)-(4*3) \rightarrow (-2)-12 \rightarrow -14$.
- b) $10\%6*8\%3$ is parenthesized as $((10\%6)*8)\%3$ since binary (multiplicative) arithmetic operators are left-associative. Evaluation: $((10\%6)*8)\%3 \rightarrow (4*8)\%3 \rightarrow (32)\%3 \rightarrow 2$.
- c) $6-3+4*5$ is parenthesized as $(6-3)+(4*5)$ since multiplicative operators bind more strongly than additive ones, and since (additive) arithmetic operators are left-associative. Evaluation: $(6-3)+(4*5) \rightarrow 3+(4*5) \rightarrow 3+20 \rightarrow 23$.
- d) $5+5*3u$ is parenthesized as $5+(5*3u)$ since multiplicative operators bind more strongly than additive ones. Evaluation: $5+(5*3u) \rightarrow 5+15u \rightarrow 20u$. Here, `int` values are implicitly converted to `unsigned int`.
- e) $31/4/2$ is parenthesized as $(31/4)/2$ since (multiplicative) arithmetic operators are left-associative. Evaluation: $(31/4)/2 \rightarrow 7/2 \rightarrow 3$.
- f) $-1-1u+1-(-1)$ is parenthesized as $(((-1)-1u)+1)-(-1)$ since (additive) arithmetic operators are left-associative. Evaluation: $(((-1)-1u)+1)-(-1) \rightarrow (((2^b-1)u-1u)+1)-(-1) \rightarrow ((2^b-2)u+1)-(-1) \rightarrow (2^b-1)u-(-1) \rightarrow (2^b-1)u-(2^b-1)u \rightarrow 0u$, due to the conversion rules.

Solution to Exercise 9. *b)* and *g)* are illegal because `-a` and `7+a`, respectively, are rvalues and the left operand of an assignment must be an lvalue. For the others, the parentheses are as follows.

- | | |
|---------------------------------|---------------------------|
| <i>a)</i> $c=((a+7)+(--b))$ | <i>c)</i> $c=(a=(-b))$ |
| <i>d)</i> $a-((a/b)*b)$ | <i>e)</i> $b*=(++a)+b$ |
| <i>h)</i> $(a+(3*(--b)))+(a++)$ | <i>f)</i> $a-(a*(+(-b)))$ |
| | <i>i)</i> $(b++)+(-a)$ |

Solution to Exercise 10. *h)* induces unspecified behavior because it is not specified whether the increment of `a` happens before or after the first `a` is evaluated. The remaining expressions can be evaluated as follows.

- a) $c=((a+7)+(--b)) \rightarrow c=((5+7)+1) \rightarrow c=(12+1) \rightarrow c=13 \rightarrow 13$.
- c) $c=(a=(-b)) \rightarrow c=(a=(-2)) \rightarrow c=(a=-2) \rightarrow c=-2 \rightarrow -2$.
- d) $a-((a/b)*b) \rightarrow 5-((5/2)*2) \rightarrow 5-(2*2) \rightarrow 5-(4) \rightarrow 1$.
- e) $b*=(++a)+b \rightarrow b*=(6+2) \rightarrow b*=8 \rightarrow 16$.

$$f) a-(a*(+(-b))) \rightarrow 5-(5*(+(-2))) \rightarrow 5-(5*(+-2)) \rightarrow 5-(5*-2) \rightarrow 5--10 \rightarrow 15.$$

$$i) (b++)+(--a) \rightarrow 2+4 \rightarrow 6.$$

Solution to Exercise 11. For any positive d , we can uniquely write n in the form

$$n = xd + y,$$

where $x, y \in \mathbb{N}$ and $y < d$. In fact, these define div and mod via $x = n \text{ div } d$ and $y = n \text{ mod } d$.

For $n = a, d = bc$, we get

$$a = pbc + q,$$

where $p, q \in \mathbb{N}$ and $q < bc$. In particular, $p = a \text{ div } (bc)$. It remains to prove that

$$(a \text{ div } b) \text{ div } c = p.$$

To this end, we use $n = q, d = b$ to get

$$q = rb + s,$$

where $r, s \in \mathbb{N}, s < b$ and $r < c$ (since $rb \leq q < cb$). Then we have

$$a = (pc + r)b + s, s < b,$$

and this means that

$$a \text{ div } b = pc + r, r < c.$$

This in turn means that $p = (a \text{ div } b) \text{ div } c$.

The implication holds as long as the mathematical value of the expression $b*c$ is representable as an unsigned `int` value. This holds since the other operations involving `/` are always error-free. If there is an overflow in $b*c$, the expressions $a/b/c$ and $a/(b*c)$ might yield different values.

Solution to Exercise 12. a) 1111 b) 10101100 c) 101001001 d) 1111111110

Solution to Exercise 13. a) 55 b) 65 c) 233 d) 341

Solution to Exercise 14.

a) 0110 b) 1100 c) 1000 d) not representable e) 1101

Solution to Exercise 15. We start by solving the Celsius-to-Fahrenheit equation for Degrees Celsius and get

$$\text{Degrees Celsius} = \frac{5(\text{Degrees Fahrenheit} - 32)}{9}.$$

The critical operation is the multiplication by 5. If \underline{m} , \overline{m} are the smallest and largest representable int values, we must guarantee that

$$\underline{m} \leq 5(\text{Degrees Fahrenheit} - 32) \leq \overline{m}$$

in order to avoid over- and underflows. Solving this for Degrees Fahrenheit yields the following bounds.

$$\frac{\underline{m}}{5} + 32 \leq \text{Degrees Fahrenheit} \leq \frac{\overline{m}}{5} + 32.$$

Since integer division rounds down for positive numerators, a valid upper bound is obtained from the expression

```
std::numeric_limits<int>::max() / 5 + 32
```

For the lower bound, we have assumed for this exercise that the integer division rounds towards zero (up) for negative numerators. Therefore, the value

```
std::numeric_limits<int>::min() / 5 + 32
```

is a valid lower bound. The resulting code is given below (the output of the mixed number simply employs the div and mod operators).

```

1 // Program: celsius.C
2 // Convert temperatures from Fahrenheit to Celsius.
3
4 #include<iostream>
5 #include<limits>
6
7 int main ()
8 {
9     // Input
10    std::cout << "Temperature in degrees Fahrenheit =? \n"
11              << " (from the interval ["
12              << std::numeric_limits<int>::min() / 5 + 32
13              << ", "
14              << std::numeric_limits<int>::max() / 5 + 32
15              << "])\n";
16    int fahrenheit;
17    std::cin >> fahrenheit;
18
19    // Computation
20    int ncelsius = 5 * (fahrenheit - 32); // numerator
21
22    // Output as mixed number
23    std::cout << fahrenheit << " degrees Fahrenheit are "
24              << ncelsius / 9 << " " << ncelsius % 9
25              << "/9 degrees Celsius.\n";
26
27    return 0;
28 }
```

Solution to Exercise 16. According to Section 2.2.8, the least significant binary digit b_0 of n is $n \bmod 2$, and the remaining digits are obtained by applying the same technique to $(n - b_0)/2 = n \operatorname{div} 2$. It follows that the second to last digit is $(n \operatorname{div} 2) \bmod 2$, and

that the third to last one is $((n \text{ div } 2) \text{ div } 2) \bmod 2$, or (as one can prove; nice Exercise for next time) $(n \text{ div } 4) \bmod 2$.

```

1 // Program: threebin.C
2 // Output the last three binary digits of a number n.
3
4 #include <iostream>
5
6 int main ()
7 {
8     // input
9     std::cout << "Last three binary digits of n =? ";
10    unsigned int n;
11    std::cin >> n;
12
13    // computation and output
14    std::cout << "The digits are "
15              << n / 2 / 2 % 2 << n / 2 % 2 << n % 2 << ".\n";
16
17    return 0;
18 }
```

Solution to Exercise 17. For part a), we first have to find a theoretical way of computing the survivor $p(k)$. One way of doing it is to generalize the problem first. Let $p_n(k)$ be the last survivor in a circle of n people when every k -th person is killed. Then $p(k) = p_{41}(k)$. Here is the crucial insight that leads to a recursive formula for $p_n(k)$:

$$p_n(k) = (p_{n-1}(k) + k) \bmod n, \quad n \geq 2,$$

where $p_1(k) = 0$ (in a circle of 1, person 0 is the only person and therefore the last survivor). To see that this formula holds, we consider the n -circle after the first person (numbered $k - 1$) has been killed. We now have an $(n - 1)$ -circle, and if we renumber the positions in such a way that the position formerly being k is now position 0, $p_{n-1}(k)$ tells us the last survivor in this new numbering. To get the survivor $p_n(k)$ in the old numbering, we simply have to add k , *but modulo* n . This takes a little thought but is easily deduced from the following table.

old numbering	0	1	...	$k - 2$	$k - 1$	k	...	$n - 1$
new numbering	$n - k$	$n - k + 1$...	$n - 2$	-	0	...	$n - k - 1$

Here is the program that determines $p(k)$ by iteratively computing the sequence $p_1(k), p_2(k), \dots, p_{41}(k)$, according to the formula above.

```

1 // Program: josephus.C
2 // determines the survivor number in the Josephus Problem
3
4 #include <iostream>
5
6 int main ()
7 {
8     // input
```

```

9  std::cout << "Every k-th person is killed for k =? ";
10 unsigned int k;
11  std::cin >> k;
12
13  // computation of  $p_{41}(k)$ 
14  unsigned int n = 1; // circle size, runs through 1, ..., 41
15  unsigned int p = 0; // survivor number, runs through  $p_1(k), \dots, p_{41}(k)$ 
16
17  // each of the following 40 identical statements increments n and updates
18  // p according to the formula  $p_n(k) = (p_{n-1}(k) + k) \bmod n$ 
19  p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n;
20  p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n;
21  p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n;
22  p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n;
23  p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n;
24  p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n;
25  p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n;
26  p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n;
27  p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n;
28  p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n; p = (p + k) % ++n;
29
30  // now  $p = p_{41}(k) = p(k)$ ; output
31  std::cout << "The survivor is the person number " << p << ".\n";
32
33  return 0;
34 }

```

Part b) is solved by playing with the program a little, and the answer is no. There are 41 possible initial positions for Josephus, and 41 ways of choosing k . Both for $k = 3$ and for $k = 7$, the survivor number is 30, so in the sequence $p(1), \dots, p(41)$, at least one of the 41 number $\{0, \dots, 40\}$ appears twice. But this implies that there must also be at least one other number in $\{0, \dots, 40\}$ that does not appear at all. The initial position with this number is therefore fatal for all $k \in \{1, \dots, 41\}$.

Interestingly, if we allow k to be *any* natural number, then it is possible to survive from any initial position.