B.6 Arrays and pointers

Solution to Exercise 77.

- a) The program outputs 5 4 1 6 0. The first array element is always the first one to be output by std::cout << *p, because p is initialized with a, and due to array-to-pointer conversion, p then points to the first element of a. The assignment p = a + *p changes p into a pointer to the element of index *p in the array (see the paragraph on adding an integer to a pointer). Since *p is initially 5, we get the element of index 5 next, which is 4. The next element is the one of index 4 (1), followed by the ones of index 1 (6) and index 6 (0). At this point, p points again to the first element of the array, so the condition p!=a fails, and the loop terminates.
- b) The general structure is this: Let i_t be the index of the array element pointed to by p after t iterations. We have $i_0 = 0$, and $i_t = a[i_{t-1}]$ for t > 0. The program terminates as soon as $i_t = 0$ again for some t > 0. We must prove that this always happens.

Assume for contradiction that 0 does not appear for a second time, so that we have an infinite loop. Since there are only n possible index values, some index distinct from 0 must appear twice in the infinite sequence i_0, i_1, \ldots Let k be some value such that $i_k \neq 0$ already appears among i_0, \ldots, i_{k-1} , and let $0 < \ell < k$ be such that $i_\ell = i_k$. By definition of the sequence, it follows that

$$\mathtt{a}[\mathfrak{i}_{k-1}] = \mathfrak{i}_k = \mathfrak{i}_\ell = \mathtt{a}[\mathfrak{i}_{\ell-1}],$$

where $k-1 \neq \ell-1$. But this gives the desired contradiction, since the array a was initialized with a sequence of pairwise distinct numbers.

Solution to Exercise 78. In a), we use

```
int* p = a+i; // address of element of index i
```

Note that

```
int* p = &a[i]; // address of element of index i
```

also works, unless i = n, since n is an out-of-bound index. A past-the-end pointer can therefore only be obtained from the first variant.

In b), we use pointer subtraction:

```
int i = p-a; // distance between *p and a[0] in the array
```

Solution to Exercise 79.

Problem 1:

The pointer b has been declared as a constant, but later it is incremented (*b++). This won't compile. Fix: remove the const. Now the program at least compiles...

Problem 2:

In the first loop, the range that p points to is too large by one address. Fix: $p \le a+7$ should be $p \le a+7$.

Problem 3

A similar problem also appears in the second loop. Inside the loop a and b are accessed by one too many values for index i. Fix: i<=7 should be i<7.

Problem 4

The cross-check in the second loop goes wrong, because b doesn't point to the beginning of the dynamically created array any more. Fix: copy the pointer b to a pointer d before the first loop, and iterate over the pointer d in the second loop. Then, we can also reintroduce the const for b.

Problem 5

The wrong type of delete-operator is used. Fix: delete should be delete[].

If all these fixed are applied, there is still one logical const missing for constcorrectness. The loop pointer p should have underlying type const int to reflect the fact that the loop does not change the array a that it loops over.

Here is the resulting correct program.

```
1
   #include < iostream >
3
   int main()
5
      int a[7] = {0, 6, 5, 3, 2, 4, 1}; // static array
     int* const b = new int[7];
6
     int*d=b;
8
a
      // copy a into b using pointers
10
     for (const int* p = a; p < a+7; ++p)
       *d++ = *p;
11
12
13
     // cross-check with random access
14
     for (int i = 0; i < 7; ++i)
15
      if (a[i] != b[i])
          std::cout << "Oops, copy error...\n";</pre>
16
17
18
     delete[] b;
19
20
      return 0;
21
   }
```

Solution to Exercise 80. We recycle the program eratosthenes2.cpp (Program 15); instead of maintaining the information whether a number has been crossed out, we maintain the information about the number of different prime divisors.

```
// Program: k_composite.cpp
// Calculate k-composite numbers in {2,...,n-1} using
// a variant of Eratosthenes' sieve.

#include <iostream>
int main()
{
```

```
// input of k
      std::cout << "Compute k-composite numbers for k =? ";</pre>
10
11
     unsigned int k;
12
     std::cin >> k;
13
14
      // input of n
     std::cout << "Compute " << k
15
16
                << "-composite numbers in \{2, \ldots, n-1\} for n =? ";
17
      unsigned int n;
18
     std::cin >> n;
19
      // definition and initialization: provides us with
20
21
      // unsigned integers composition [0], \ldots, composition [n-1]
22
      unsigned int* const composition = new unsigned int [n];
      for (unsigned int i = 0; i < n; ++i)
23
24
       composition[i] = 0; // no information yet
25
26
      // computation and output
27
     std::cout << k << "-composite numbers in {2,...,"
                << n-1 << "}:\n";
28
      for (unsigned int i = 2; i < n; ++i) {
29
       if (composition[i] == 0) {
30
31
         /\!/ i is prime: add 1 to composition number of all
32
          // multiples (including i)
33
         for (unsigned int m = i; m < n; m += i)
34
            ++composition[m];
35
       // now the composition number of i is up-to-date
36
37
        if (composition[i] == k)
38
         std:: cout << i << " ";
39
40
     std::cout << "\n";
41
                                               // free dynamic memory
42
     delete[] composition;
43
44
     return 0:
45
   }
```

The following are the 8 different 7-composite numbers smaller than 1,000,000: 510510, 570570, 690690, 746130, 870870, 881790, 903210, 930930.

Solution to Exercise 81. One trick here is that computations with indices modulo 3 save us the signs $(-1)^{i+j}$.

```
// Prog: inverse_matrix.cpp
   // read in a 3x3 matrix A, compute ins inverse A^{-1},
   // and output it along with A x A^{-1} as a crosscheck
   #include < iostream >
6
   int main()
7
8
9
     // read in A (as a sequence of 9 numbers)
     double a[3][3];
10
     for (int i=0; i<3; ++i)
11
      for (int j=0; j<3; ++j)
12
13
         std::cin >> a[i][j];
14
15
     // compute determinant of A via Sarrus' rule
     double det = 0;
16
     for (int i=0; i<3; ++i)
17
       det += a[0][i] * a[1][(i+1)%3] * a[2][(i+2)%3]
18
```

```
19
          - a[2][i] * a[1][(i+1)%3] * a[0][(i+2)%3];
20
      // compute (and output) entries of A^{-1} through Cramer's rule
21
      std::cout << "A^-1 = \n";
22
23
      double a_inv[3][3];
24
      for (int i=0; i<3; ++i) {
25
        for (int j=0; j<3; ++j)
26
          std::cout
27
            << (
                 a_{inv}[i][j] = ((a[(j+1)%3][(i+1)%3]*a[(j+2)%3][(i+2)%3]) -
28
                                  (a[(j+1)%3][(i+2)%3]*a[(j+2)%3][(i+1)%3])) / det)
29
30
            << " ";
31
        std::cout << "\n";
32
33
34
      // crosscheck
35
      std::cout << "A * A^-1 = \n";
      for (int i=0; i<3; ++i) {
36
37
       for (int j=0; j<3; ++j)
          // output (A x A^{-1})_{ij}
38
39
          std::cout <<
            a[i][0]*a_inv[0][j]+
40
41
            a[i][1]*a_inv[1][j]+
        a[i][2]*a_inv[2][j] << " ";
std::cout << "\n";
42
43
44
45
46
      return 0;
47
48
   }
```

Solution to Exercise 82.

```
// Program: read_array.cpp
   // read a sequence of n numbers into an array
3
   #include <iostream>
5
   int main()
6
   {
     // input of n
     unsigned int n;
8
9
     std::cin >> n;
10
      // dynamically allocate array
11
12
     int* const a = new int[n];
13
14
      // read into the array
     for (int i=0; i<n; ++i) std::cin >> a[i];
15
16
17
      // output what we have
     for (int i=0; i<n; ++i) std::cout << a[i] << " ";
18
     std::cout << "\n";
19
20
     // delete array
21
22
     delete[] a;
23
24
     return 0;
25
26
```

Solution to Exercise 83.

```
1 // Program: sort_array.cpp
   // read a sequence of n numbers into an array,
// sort them, and output the sorted sequence
   #include <iostream>
6
   int main()
7
      // input of n
8
      unsigned int n;
a
      std::cin >> n;
10
11
12
      // dynamically allocate array
13
      int* const a = new int[n];
14
15
      // read into the array
16
      for (int i=0; i<n; ++i) std::cin >> a[i];
17
      // sort array: in round i=0,...,n-2 we find
18
      // the smallest element in a[i], \ldots, a[n-1] and
19
20
      // interchange it with a[i]
      for (int i=0; i<n-1; ++i) {
21
        // find minimum in a[i], \ldots, a[n-1]
22
23
        int i_min = n-1; // index of minimum
24
        for (int j=i; j < n-1; ++j)
          if (a[j] < a[i_min]) i_min = j;</pre>
25
        // interchange a[i] with a[i_min]
26
27
        const int h = a[i]; a[i] = a[i_min]; a[i_min] = h;
28
29
30
      // output sorted sequence
31
      for (int i=0; i<n; ++i) std::cout << a[i] << " ";
      std::cout << "\n";
32
33
34
      // delete array
35
      delete[] a;
36
37
      return 0;
38
39
   }
```

Solution to Exercise 84.

```
1 // Program: cycles.cpp
   // read a sequence of n numbers into an array; if the sequence // encodes a permutation of \{0,\ldots,n-1\}, output its cycle
    // decomposition
   #include <iostream>
 8
    int main()
9
    {
10
      // input of n
11
      unsigned int n;
12
       std::cin >> n;
13
14
       \ensuremath{/\!/} dynamically allocate array for the numbers, and a
15
       /\!/\ second\ array\ of\ booleans\ to\ keep\ track\ of\ which
16
      // numbers are present
17
      int* const a = new int[n];
18
      bool* const present = new bool[n];
19
20
      // initialize array present
```

```
21
      for (bool* r=present; r<present+n; ++r) *r = false;</pre>
22
23
      // read number into the array and remember that it was read
24
      for (int* p=a; p < a+n; ++p) {
25
        std::cin >> *p;
        if (*p >= 0 && *p < n)
26
27
          present[*p] = true;
28
29
30
      // check whether we have read all numbers in \{0, \ldots, n-1\}
31
      bool ok = true;
      for (bool* r=present; r<present+n; ++r)</pre>
32
33
        if (!*r) {
34
          std::cout << "input sequence does not encode a permutation.\n";</pre>
35
          ok = false;
36
          break;
37
38
      if (ok) {
39
40
        // do the cycle decomposition. Here we reuse the array present
41
        // and remove from it all numbers that we have already put into
        // some cycle
42
        std::cout << "cycle decomposition is ";</pre>
43
44
        int next = 0; // next number not yet put into a cycle
        while (next < n) {
45
46
          // output cycle starting with next; we must come back to next
47
          // below: assuming we would come back to some other element
          // on the cycle, that element would have two preimages under
48
49
          // pi, a contradiction
50
          const int first = next;
          std::cout << "( ";
51
52
          do {
           std::cout << next << " ";
53
54
            present[next] = false;
            next = a[next]; // next -> pi(next)
          } while (next != first);
56
57
          std::cout << ") ";
58
59
          // find start element of next cycle
60
          while (!present[next]) ++next;
61
62
        std::cout << "\n";
63
      // delete arrays
64
65
      delete[] present;
66
      delete[] a;
67
68
      return 0;
69
70
```

Solution to Exercise 85. Let $s=\alpha\ldots ab$ (i.e. m-1 a's followed by one b); let $t=\alpha\ldots a$ (i.e. n a's). Then the algorithm must always go through all m characters of any window in order to find the mismatch with b at the last position. Since it in total processes n-m+1 windows $\{1,\ldots,m\}$ up to $\{n-m+1,\ldots,n\}$, the number of comparisons is m(n-m+1).

Solution to Exercise 86.

```
1 // Program: threedim_array.cpp
```

```
// iterate over a multidimensional array
    #include <iostream>
    int main()
 5
 6
       int a[4][2][3] =
          { \{ \ // \ the \ 4 \ elements \ of \ a: }
 8
 9
             { // the 2 elements of a[0]:
                \{2, 4, 5\}, //  the three elements of a[0][0]
10
                \{4, 6, 7\} // the three elements of a [0][1]
11
12
             },
             { // the 2 elements of a[1]:
13
                {1, 5, 9}, // the three elements of a[1][0] {4, 6, 1} // the three elements of a[1][1]
14
15
16
17
             { // the 2 elements of a[2]:
                \{5, 9, 0\}, // \text{ the three elements of a [2] [0]} \{1, 5, 3\} // \text{ the three elements of a [2] [1]}
18
19
20
21
             \{ // \text{ the 2 elements of a[3]:} 
                {6, 7, 7}, // the three elements of a[3][0]
{7, 8, 5} // the three elements of a[3][1]
22
23
             }
24
25
          };
26
       for (const int (*i)[2][3] = a; i < a + 4; ++i) {
27
28
          // i (pointer to int[2][3]) points to a[0],...,a[3]
          //*i therefore assumes the values a[0],...,a[3]
29
30
          for (const int (*j)[3] = *i; j < *i + 2; ++j) {
            // j (pointer to int[3]) points to a[i][0],...,a[i][1]
// *j therefore assumes the values a[i][0],...,a[i][1]
31
32
33
             for (const int* k = *j; k < *j + 3; ++k)
                // k (pointer to int) points to a[i][j][0], \dots a[i][j][2] // *k therefore assumes the values a[i][j][0], \dots a[i][j][2]
34
35
                std::cout << *k << " ";
37
             std::cout << "\n";
38
39
          std::cout << "\n";
       }
40
41
42
       return 0;
43
```

Some people might be tempted by the following kind of approach:

```
1
   // Program: threedim_array.cpp
   // (erroneously) iterate over a multidimensional array
   #include <iostream>
5
   int main()
6
   {
7
     int a[4][2][3] =
       8
q
            {2, 4, 5}, // the three elements of a[0][0]
10
11
            {4, 6, 7} // the three elements of a[0][1]
12
         },
13
         \{ // \text{ the 2 elements of a[1]:} 
14
            \{1, 5, 9\}, // the three elements of a[1][0]
15
           \{4, 6, 1\} // the three elements of a[1][1]
16
         \{ \ // \ the \ 2 \ elements \ of \ a[2]:
17
```

```
\{5, 9, 0\}, // the three elements of a[2][0]
18
19
               \{1, 5, 3\} // the three elements of a[2][1]
21
            { // the 2 elements of a[3]:
              \{6, 7, 7\}, // \text{ the three elements of a } [3][0]
\{7, 8, 5\} // \text{ the three elements of a } [3][1]
22
23
            }
24
25
26
27
       const int* p = a[0][0];
                                              // pointer to a[0][0][0]
       for (int i=0; i<24; ++i)
28
         std::cout << *p++ << " ";
29
30
31
       std::cout << "\n";
32
33
       return 0;
34
```

This indeed does not contradict anything written in Section 2.6; in particular, any operation ++p (which reduces to p+1 plus an assignment) has the property that both p as well as p+1 point to elements (or past the end) of the *same* array. However, that array changes during the increment. For the first three ++p's, it's the array a[0][0], for the second three, it's a[0][1], and so on. But changing the array during pointer increment is not allowed by the C++standard; in fact, the standard allows (by not forbidding it) implementations of pointer arithmetic that perform bounds checking. Such an implementation might give you a runtime error if you try to increment p further than past the end of a[0][0], the array on which the pointer logically "lives".

Solution to Exercise 87. The trick is to use characters (which have integral values) directly as array indices.

```
// Program: frequencies.cpp
   // output frequencies of the letters in an input text
  #include < iostream >
5
6
   int main ()
      //\ array\ for\ number\ of\ occurences\ of\ every\ ASCII\ character
8
9
      int frequency[128];
     for (int i=0; i<128; ++i) frequency[i] = 0;
10
11
12
      // now scan the text
                                 // next character
13
     char c:
14
      unsigned int total = 0;
                                 // text length
15
     while (std::cin >> c) {
16
       ++total:
17
        ++frequency[c];
18
19
     // output
20
21
     unsigned int letters = 0; // number of letters
      std::cout << "Frequencies: \n";</pre>
22
      for (char c = 'a'; c <= 'z'; ++c) {
24
       const int f = frequency[c] + frequency[c-32]; // lower + upper case c
25
        letters += f;
       std::cout << c << ": " << f << " of " << total << "\n";
26
27
```

```
28 std::cout << "Other: " << total-letters << " of " << total << "\n";
29
30 return 0;
31 }
```

Solution to Exercise 88. Here is a solution. We initially dynamically allocate an array of length n=1 (pointed to by a pointer a), and whenever the next sequence element wouldn't fit anymore, we replace the array by a new one of length 2n. For this, we first dynamically allocate a helper array, copy the contents of the current array into the helper array, delete the current array and then let a point to the newly allocated helper array.

```
// Program: read_array.cpp
   // read a sequence of numbers into an array
   #include <iostream>
3
   int main()
6
   {
7
     int n = 1; // current array size
     int k = 0; // number of elements read so far
8
q
10
     // dynamically allocate array
     int* a = new int[n]; // this time, a is NOT a constant
11
12
13
     // read into the array
14
     while (std::cin >> a[k]) {
15
       if (++k == n) {
16
         // next element wouldn't fit; replace the array a by
         // a new one of twice the size
17
         int* b = new int[n*=2]; // get pointer to new array
18
19
         for (int i=0; i<k; ++i) // copy old array to new one
20
           b[i] = a[i];
21
         delete[] a;
                                    // delete old array
                                    // let a point to new array
22
         a = b;
23
       }
24
25
26
     // output the first k elements
27
     for (int i=0; i<k; ++i) std::cout << a[i] << " ";
     std::cout << "\n";
28
29
     // delete array
30
31
     delete[] a;
32
33
     return 0:
34
35
```

This is space and time efficient. The constant of proportionality in (i) is 3: whenever we grow the array (and these are the points in time where the ratio between memory cells in use and k is largest), we allocate a new array of length 2n = 2k, in addition to the one of length k that we already have. This means that we have 3k memory cells in use at that time.

The constant of proportionality in (ii) is 3 as well. To see this, let us consider the situation after an execution of the while loop, where k is the number of elements read so far. There have been k assignments to array elements in the loop's condition, and some

additional assignments during the copying of old to new array. Such assignments took place when the number of elements currently read was a power of two less or equal to k, and the number of these additional assignments was exactly the power of two in question. Since the sum of all powers of two less or equal to k is at most $k + k/2 + k/4 + \cdots \le 2k$, the total number of assignments is bounded by k + 2k = 3k.

Solution to Exercise 89. Here is the faster program.

```
1 #include < iostream >
2 #include < cassert >
   int main()
5
      // read floor dimensions
      int n; std::cin >> n; // number of rows
      int m; std::cin >> m; // number of columns
8
      // dynamically allocate twodimensional array of dimensions
10
      // (n+2) x (m+2) to hold the floor plus extra walls around
11
12
      int** const floor = new int*[n+2];
13
      for (int r=0; r < n+2; ++r)
14
        floor[r] = new int[m+2];
15
16
      // we need another two arrays for storing row and column
      // indices of already labeled cells;
17
      int* const labeled_r = new int[n*m];
18
      int* const labeled_c = new int[n*m];
19
20
21
      // in order to search for new cells to be labeled, we
22
      // always start from the first labeled cell whose neighbors
      // have not been looked at yet;
23
24
      int next_1 = 0; // index of this cell
25
      // whenever we label a cell, we append it to the list of
26
27
      // labeled cells; that way, the cells are ordered by label
28
      // in the list
      int last_1 = 0; // one plus index of last cell in this list
29
30
      // target coordinates, set upon reading 'T'
31
32
      int tr = 0;
      int tc = 0;
33
34
35
      // assign initial floor values from input:
      // source: 'S' -> 0 (source reached in 0 steps)
// target: 'T' -> -1 (number of steps still unkn
// wall: 'X' -> -2
// empty cell: '-' -> -1 (number of steps still unkn
36
37
                                   -1 (number of steps still unknown)
38
                                   -1 (number of steps still unknown)
39
      for (int r=1; r<n+1; ++r)
40
41
        for (int c=1; c<m+1; ++c) {
          char entry = '-';
42
43
          std::cin >> entry;
                  (entry == 'S') {
44
          if
            floor[r][c] = 0;
45
46
            labeled_r[last_l] = r;
47
            labeled_c[last_l] = c;
48
             ++last_l;
49
          }
          else if (entry == 'T') floor[tr = r][tc = c] = -1;
50
51
          else if (entry == 'X') floor[r][c] = -2;
          else if (entry == '-') floor[r][c] = -1;
52
53
```

```
54
55
       // add surrounding walls
      for (int r=0; r<n+2; ++r)
56
57
        floor[r][0] = floor[r][m+1] = -2;
58
      for (int c=0; c < m+2; ++c)
59
        floor[0][c] = floor[n+1][c] = -2;
60
61
      // main loop: process next labeled cell until done
62
      while (next_l != last_l) {
        const int r = labeled_r[next_l];
63
64
        const int c = labeled_c[next_l];
        const int i = floor[r][c];
65
66
        assert (i >= 0);
67
        // label the unlabeled neighbors by i+1
        for (int rr = r-1; rr <= r+1; ++rr)
68
69
          for (int cc = c-1; cc <= c+1; ++cc)
70
            if ( (rr == r \mid \mid cc == c) \&\& floor[rr][cc] == -1) {
71
               // we have a neighbor, and it's not labeled yet
              floor[rr][cc] = i+1;
72
73
               labeled_r[last_1] = rr;
74
               labeled_c[last_1] = cc;
75
               ++last_l;
            }
76
77
        ++next_1;
78
79
80
      // mark shortest path from source to target (if there is one)
      int r = tr; int c = tc; // start from target
81
82
      while (floor[r][c] > 0) {
83
        const int d = floor[r][c] - 1; // distance one less
        floor[r][c] = -3; // mark cell as being on shortest path
84
85
        // go to some neighbor with distance d
86
                (floor[r-1][c] == d) --r;
        if
        else if (floor[r+1][c] == d) ++r;
87
        else if (floor[r][c-1] == d) --c;
                                      ++c; // (floor[r][c+1] == d)
89
        else
90
91
      // print floor with shortest path
92
93
      for (int r=1; r<n+1; ++r) {
94
        for (int c=1; c<m+1; ++c)
95
          i f
                  (floor[r][c] == 0)
                                       std::cout << 'S';
96
          else if (r == tr && c == tc) std::cout << 'T';
          else if (floor[r][c] == -3) std::cout << 'o';
97
98
          99
          else
                                        std::cout << '-';
        std::cout << "\n";
100
101
102
103
      // delete dynamically allocated arrays
      delete[] labeled_c;
104
105
      delete[] labeled_r;
106
      for (int r=0; r< n+2; ++r)
107
        delete[] floor[r];
108
      delete[] floor;
109
110
      return 0:
111 }
```

Solution to Exercise 90. The idea is to try all possible combinations of a, b, c within prespecified ranges. In order to make this fast, some tricks are needed, though. First of all, we use Eratosthenes' Sieve in order to precompute the information whether a given

number that may arise as $|an^2 + bn + c|$ is prime. Then we use another array to mark the primes that we have seen in a run of primes; the trick here is to delete the markers again in an efficient way. The following program discovers a quadratic polynomial of Euler quality 45, namely

```
36n^2 - 810n + 2753.
```

```
// Program: euler_prime.cpp
  // Finds a,b,c within specified bounds such that
3 // the formula |an^2 + bn + c| produces the largest
   // number of distinct consecutive primes, starting
   // with n = 0
   #include <iostream>
   #include <cassert>
8
10
   int main ()
11 {
      // by multiplying with -1, if necessary, we may assume c \ge 0
12
      const int arange = 100;
                                   // a in {-arange+1,...,arange-1}
13
                                   // b in {-brange+1,...,brange-1}
      const int brange = 1000;
14
                                   // c in {2,..., crange-1}
15
      const int crange = 10000;
16
17
      // first, compute all primes in the set \{an^2 + bn + c\} where a, b, c
      // run through their ranges and n is smaller than 100 (we're not
18
      // searching for longer runs of primes here)
19
20
      const int max_elem = arange * 10000 + brange * 100 + crange;
21
22
      // run Eratosthenes' sieve
23
      bool prime[max_elem];
      for (int i = 2; i < max_elem; ++i)</pre>
24
25
       prime[i] = true;
      for (int i = 2; i < max_elem; ++i)</pre>
26
27
       if (prime[i]) {
28
          // cross out all proper multiples of i
29
          for (int m = 2*i; m < max_elem; m += i)
            prime[m] = false;
30
31
32
33
      // now search for the best a, b, c
      int best_run = 0;
34
35
      int best_a = arange;
36
      int best_b = brange;
      int best_c = crange;
37
38
39
      // array to keep track of primes we have already seen in a run
40
      bool seen before [max elem]:
41
      for (int i = 2; i < max_elem; ++i)</pre>
42
        seen_before[i] = false;
43
44
      // loop over all candidate triples (a,b,c)
      for (int c = 2; c < crange; ++c) {</pre>
45
        if (!prime[c]) continue; // not even a prime for n = 0
46
47
        for (int a = -arange+1; a < arange; ++a)</pre>
          for (int b = -brange+1; b < brange; ++b) {
48
            // evaluate elem = an^2 + bn + c for n=0,1,...
49
            int n = 0;
51
            int elem = c;
52
            for (; n < 100; ++n) {
53
             const int abs_elem = elem < 0 ? -elem: elem;</pre>
              if (abs_elem < 2 || !prime[abs_elem]) break; // not a prime
```

```
55
               if (seen_before[abs_elem]) break;
                                                                // repeated prime
56
               seen_before[abs_elem] = true;
                                                                 // new prime
               // update element
58
              elem += a * (2*n + 1) + b;
59
60
            // now we have seen a run of n primes (for 0, \ldots, n-1)
61
            if (n > best_run) {
62
              best_run = n;
63
              best_a = a;
              best_b = b;
64
            best_c = b;
best_c = c;
}
65
66
67
            // remove the "seen_before" markers for the next run
68
             elem = c;
            for (int k=0; k<n; ++k) {
69
70
              const int abs_elem = elem < 0 ? -elem: elem;</pre>
71
               seen_before[abs_elem] = false;
72
               elem += a * (2*k + 1) + b;
73
74
75
     }
76
      std::cout << "Best a
                                    = " << best_a << ".\n";
77
                                    = " << best_b << ".\n";
78
      std::cout << "Best b
      std::cout << "Best b = " << best_b << ".\n";
std::cout << "Best c = " << best_c << ".\n";
79
     std::cout << "Euler quality = " << best_run << ".\n";
ጸበ
81
82
     return 0:
83 }
```

Solution to Exercise 91.

```
1 // Prog: xbm.cpp
2 // includes an xbm file and outputs the xbm-file that corresponds
    // to the image rotated by 90 degrees
   #include<iostream>
 6 #include < cassert >
   #include "original.xbm"
8 #define width original_width
9 #define height original_height
10 #define bits original_bits
11
12 int main()
13
      // array for mapping from {0,...,15} to {'0',...,'f'} char hex[] = {'0', '1', '2', '3', '4', '5', '6', '7', '8', '9', 'a', 'b', 'c', 'd', 'e', 'f'};
14
15
16
17
18
       // output header for rotated image
       std::cout << "#define rotated_width " << height << "\n";
std::cout << "#define rotated_height " << width << "\n";</pre>
19
20
21
       std::cout << "static unsigned char rotated_bits[] = {\n";</pre>
22
23
       // go through the pixels columnwise (from right to left),
       \ensuremath{//} and within each column, proceed from top to bottom
24
       const unsigned int columns = (width+7)/8; // bytes in row (original) const unsigned int rows = (height+7)/8; // bytes in column (rotated)
25
26
                                                    // not before first byte
27
       bool comma = false;
28
       for (int c=columns-1; c>=0; --c) {
29
         // go through columns from right to left
         for (int d=128; d>0; d/=2) {
30
31
           // for d = 2^j, we are in column c+j
```

```
for (int r=0; r< rows; ++r) {
32
               // go through rows from top to bottom and build
// up one byte from any 8 pixels (least significant
// digit comes from pixel in first of these rows)
33
35
36
               int byte = 0;
37
               for (int i=7; i>=0; --i) {
38
                 // we are in row 8r+i
39
                  int pixel;
                 if (8*r+i >= height)
40
                    // non-existing row, fill up with zeros
41
                    pixel = 0;
42
43
                  else
44
                    // get pixel in row 8r+i and column c+j
45
                    pixel = (unsigned char)(bits[(8*r+i)*columns+c])/d%2;
                 byte = 2*byte+pixel;
46
47
               if (comma) std::cout << ", "; comma = true;
std::cout << "0x" << hex[byte/16] << hex[byte%16];</pre>
48
49
50
51
          }
52
          std::cout <<"\n";
53
54
       std::cout << "}\n;";
55
56
       return 0;
57
58
    }
```