

## Approximation Algorithms and Semidefinite Programming FS12 Exercise Set 6

Course Webpage: <http://www.ti.inf.ethz.ch/ew/lehre/ApproxSDP12/>

Discussion: Mai 22 & 29

### Exercise 1 (Vector discrepancy)

[Exercise 11.1]

- Show that every set system on  $n$  points has vector discrepancy at most  $\sqrt{n}$ .
- Show that this bound is tight, possibly up to a multiplicative constant independent of  $n$ .
- Let  $\mathcal{F}$  be a system on  $n$  sets on a set  $V$  of  $n$  points, such that every point is contained in exactly  $r$  sets of  $\mathcal{F}$ , and for every two distinct points  $i, j \in V$ , there are exactly  $t$  sets  $F \in \mathcal{F}$ , with  $\{i, j\} \subset F$ . Prove that  $\text{vecdisc}(\mathcal{F}) \geq \sqrt{r-t}$ .

**Hint - Exercise 1.c)** Instead of  $\max_{F \in \mathcal{F}} \|\sum_{j \in F} \mathbf{u}_j\|^2$ , estimate  $\sum_{F \in \mathcal{F}} \|\sum_{j \in F} \mathbf{u}_j\|^2$ .

### Exercise 2 (Discrepancy I)

[Exercise 11.2] Use the probabilistic method to show that  $\text{disc}(\mathcal{F}) = O(\sqrt{n \log m})$  for every system of  $m$  sets on  $n$  points.

**Hint - Chernoff bound:** Let  $X_1, \dots, X_n$  be independent random variables with  $0 \leq X_i \leq 1$ , for  $i = 1, \dots, n$  and let  $X = X_1 + \dots + X_n$ ,  $\mu = \mathbb{E}[X]$ . Then for all  $0 \leq \epsilon \leq 1$ :  $\Pr[|X - \mu| \geq \epsilon\mu] \leq 2 \exp\left(-\frac{\epsilon^2}{3}\mu\right)$ .

### Exercise 3\* (Discrepancy II) [\*Bonus exercise]

[Exercise 11.3] Use the probabilistic method to show the existence of set systems with  $n^2$  sets on  $n$  points and with discrepancy  $\Omega(\sqrt{n \log n})$ .

(Together with Exercise 1.a) this shows that the gap between  $\text{vecdisc}$  and  $\text{disc}$  can be at least of order  $\sqrt{\log m}$ , for  $m = n^2$ . The complete set system  $2^V$  exhibits a similar gap for  $m = 2^n$ .)

**Hint:** Fix an arbitrary coloring  $\chi$  and show that the discrepancy of  $\chi$  for a system of  $n^2$  independent random sets is below  $c\sqrt{n \log n}$  with probability smaller than  $2^{-n}$ . A random set is obtained by including each point independently with probability  $\frac{1}{2}$ .